

Discussion on:
Relative Price Dispersion: Evidence and Theory
by Kaplan, Menzio, Rudanko and Trachter

Gaetano GABALLO

Banque de France

17 December 2015

The paper

- ▶ **Relative Price Dispersion:** persistent difference in prices for the same good across shops, which is not explained by shop-specific average price
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- ▶ Evidence: impressive dataset allows for detailed price decomposition
- ▶ Theory: extend a canonical model to account for relative price dispersion
 - ▶ model: Burdett and Judd (1983)
 - ▶ addition: shops sell *two* goods and try to discriminate *heterogeneous* buyers

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$$\Pi(\bar{p}) = \{\alpha + 2(1 - \alpha)[1 - G(p)]\}p$$

where \bar{p} is the reservation price, $\Pi(\bar{p}) \equiv \alpha\bar{p}$ and $G(p)$ is the fraction of sellers with a price less than p .

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- ▶ If p is the av. price of a basket of goods, then RPD (indet.) even with $\alpha = 1$

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- ▶ To understand the importance of each assumption, let us see how far I can go without relying on the BJ model.

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- C If cool cannot buy the good then go to the other shop

Simpler Game: strategies

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crucial non-linearity: Given ass. H, \mathcal{B} is strictly dominated by \mathcal{D} !

Payoff table

vs \rightarrow	C	$D1$	$D2$
C	$2p_c$	$2.5p_c$	$2.5p_c$
$D1$	$0.5p_c + p_b$	$0.5p_c + p_b$	$p_c + p_b$
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With $p_b > p_c > \frac{2}{3}p_b$ only (C, C) is Nash.

Given ass. C , price dispersion has private (and **social!**) value!

Conclusion

- ▶ Important evidence on the persistence of price dispersion
- ▶ The new assumptions alone already buy the existence of equilibria with RPD
- ▶ BJ introduces indeterminacy without giving clear advantages
- ▶ **Beautiful and challenging agenda!**
Huge literature on price discrimination in IO: how to properly incorporate these insights into the Macro literature?